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PATENT TRADEMARK OFFICE

This application is submitted in the name of the following inventors:

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The assignee is Network Appliance, Inc., a corporation having an office at
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Title of the Invention

File Access Control in a Multi-protocol File Server

Background of the Invention

1. Field of the Invention

The invention relates to file access control in a multi-protocol file server.

2. *Related Art*

In an integrated computer network, it is desirable for multiple client devices to share access to the same files. One known method is to provide a network file server for storing files, capable of receiving and responding to file server requests from those client devices. These file server requests are made using a file server protocol, which is recognized and adhered to by both the file server and the client device. Because the files are stored at the file server, multiple client devices have the opportunity to share access to the same files.

In a file system intended for use by more than one user, it is desirable to restrict access by programs to files in the file system. Restricting access includes at least the aspects of (1) user authentication—determining that requesting users are truly who they say they are, and (2) access control validation—determining that an authenticated user is allowed to access a particular file in a particular way. When the file system is maintained on a file server remote from the user making the request, there is an additional aspect of the access control protocol—what requests can be made by the user to access files or to set access control for files.

One problem in the known art is that there are multiple diverse models for access control validation, each typically associated with a particular file system, and there are multiple diverse access control protocols, each typically corresponding to a model for

1 access control validation. Despite the differences between these models and protocols,
2 the file server should respond to file server requests from each user, and should exhibit
3 access control validation behavior, consistent with each user's model and without secu-
4 rity violations or surprises for users.

5
6 For example, a first access control model in common use is associated with
7 the Unix operating system (or a variant thereof). This first access control model associ-
8 ates permissions with each file for a file owner, an owner's group, and all other users.
9 These permissions allow access (for the owner, group, or all other users) to read, write, or
10 execute the indicated file. This first access control model is typically implemented by the
11 NFS ("Network File System") file server protocol, possibly augmented with an adjunct
12 file-locking protocol, NLM ("Network Lock Manager"). A second access control model
13 in common use is associated with the Windows NT operating system. This second access
14 control model associates an ACL (access control list) with each file, each entry in the
15 ACL specifying an individual user, a group of users, or all users. Each entry can allow
16 access (for the specified users) to read, write, or execute the indicated file, or can specifi-
17 cally deny access. This second access control model is typically implemented by the
18 CIFS ("Common Internet File System") protocol. However, NT devices can also use the
19 NFS protocol by means of the "PC NFS" implementation, and Unix devices can also ma-
20 nipulate POSIX ACLs. These two access control models in common use differ in sig-
21 nificant ways, including (1) what permissions can be assigned to a file, (2) with what

1 granularity of specificity permissions can be assigned, and (3) how users are identified so
2 as to match them with permissions.

3
4 One method known in the art is to provide a multi-protocol file server that
5 maps all security semantics to that of a single native operating system for the file server,
6 and uses that single native operating system to validate file access control. The "Samba"
7 system and similar emulation packages are believed to use this known method. This
8 known method has the drawback that it can result in security errors or surprises for those
9 client devices using security semantics other than the file server's native operating sys-
10 tem.

11
12 Another method known in the art is to provide a multi-protocol file server
13 that supports differing types of security semantics for differing files, but attempts to vali-
14 date file access control for each user using the user's access control model. Some "Net-
15 ware" products available from Novell Corporation are believed to use this known
16 method. This known method has the drawback that the user's access control model can
17 differ significantly from the access control model set for the file, resulting in security er-
18 rors or surprises for those client devices using security semantics other than associated
19 with the target file.

20
21 Accordingly, it would be desirable to provide a method and system for en-
22 forcing file security semantics among client devices using multiple diverse access control

1 models and multiple diverse file server protocols. This advantage is achieved in an em-
2 bodiment of the invention in which a multi-protocol file server identifies each file with
3 one particular access control model out of a plurality of possible access control models,
4 and enforces that particular access control model for all accesses to that file. When the
5 file server receives a file server request for that file using a file server protocol with a dif-
6 ferent access control model, the file server translates the access control limits imposed by
7 the file's access control model into no-less-restrictive access control limits in the different
8 access control model. The file server restricts access to the file using the translated ac-
9 cess control limits.

10 11 Summary of the Invention

12
13 The invention provides a method and system for enforcing file access con-
14 trol among client devices using multiple diverse access control models and multiple di-
15 verse file server protocols. A multi-protocol file server identifies each file with one par-
16 ticular access control model out of a plurality of possible models, and enforces that one
17 particular model for all accesses to that file. When the file server receives a file server
18 request for that file using a different access control model, the file server translates the
19 access control limits for that file into no-less-restrictive limits in the different model. The
20 file server restricts access by the client device using the translated access control limits.

1 In a preferred embodiment, each file is assigned the access control model of
2 the user who created the file or who last set access control limits for the file. When a user
3 having a different access control model sets access control limits, the access control
4 model for the file is changed to the new model. Files are organized in a tree hierarchy, in
5 which each tree is limited to one or more access control models (which can limit the abil-
6 ity of users to set access control limits for files in that tree). Each tree can be limited to
7 NT-model-only format, Unix-model-only format, or mixed NT-or-Unix-models format.

8 9 Brief Description of the Drawings

10
11 The figure shows a block diagram of a system for enforcing diverse access
12 control models among client devices.

13 14 Detailed Description of the Preferred Embodiment

15
16 In the following description, a preferred embodiment of the invention is de-
17 scribed with regard to preferred process steps and data structures. However, those skilled
18 in the art would recognize, after perusal of this application, that embodiments of the in-
19 vention may be implemented using one or more general purpose processors (or special
20 purpose processors adapted to the particular process steps and data structures) operating
21 under program control, and that implementation of the preferred process steps and data

1 structures described herein using such equipment would not require undue experimenta-
2 tion or further invention.

3
4 Inventions described herein can be used in conjunction with inventions de-
5 scribed in the following applications:

6
7 o Application Serial No. 08/985,398, filed December 5, 1997, in the name of Andrea
8 Borr, titled "Multi-Protocol Unified File-locking", attorney docket number NET-
9 023.

10
11 This application is hereby incorporated by reference as if fully set forth
12 herein.

13
14 The enclosed technical appendix is part of the disclosure of the invention,
15 and is hereby incorporated by reference as if fully set forth herein.

16
17 *System Elements*

18
19 The figure shows a block diagram of a system for enforcing diverse access
20 control models among client devices.

21
22 A system 100 includes a file server 110, and a set of client devices 120.

1
2 The file server 110 maintains a file system 111, including a set of files 112.

3
4 In a preferred embodiment, the file server 110 maintains the file system 111
5 using inventions described in the following patent applications:
6

7 o Application Serial No. 08/471,218, filed June 5, 1995, in the name of inventors
8 David Hitz et al., titled "A Method for Providing Parity in a Raid Sub-System
9 Using Non-Volatile Memory", attorney docket number NET-004;

10
11 o Application Serial No. 08/454,921, filed May 31, 1995, in the name of inventors
12 David Hitz et al., titled "Write Anywhere File-System Layout", attorney docket
13 number NET-005;

14
15 o Application Serial No. 08/464,591, filed May 31, 1995, in the name of inventors
16 David Hitz et al., titled "Method for Allocating Files in a File System Integrated
17 with a Raid Disk Sub-System", attorney docket number NET-006.

18
19 Each of these applications is hereby incorporated by reference as if fully set
20 forth herein.
21

1 The file server 110 is disposed for receiving file server requests 121 from
2 the client devices 120. The file server 110 parses each request 121, determines whether
3 the operation requested in the request 121 is allowed (for the client device 120 that sent
4 the request 121 and for the one or more target files 112 specified by the request 121). If
5 allowed, the file server 110 performs that operation on the one or more target files 112.

6
7 The file server 110 is also disposed for transmitting file server responses
8 122 to the client devices 120. The file server 110 determines the response to each request
9 121 (possibly including a response indicating that the requested operation was not al-
10 lowed), generates that response 122, and transmits that response 122 to the client device
11 120 that sent the request 121.

12
13 Each client device 120 is disposed for transmitting file server requests 121
14 to the file server 110, and for receiving file server responses 122 from the file server 110.

15
16 *Access Control Models*

17
18 In a preferred embodiment, each client device 120 can be either a Unix cli-
19 ent device 120 or a Windows NT client device 120. Each client device 120 can either use
20 the NFS file server protocol to make requests 121, or use the CIFS file server protocol to
21 make requests 121. (Although typically Unix client devices 120 use the NFS file server
22 protocol and NT client devices 120 use the CIFS file server protocol, it is possible for NT

client devices 120 to use the NFS file server protocol by using the PC-NFS implementation of that file server protocol.) The file server 110 receives each request 121 and (if allowed) performs the requested operation on the target files 112 specified by the request 121.

The file server 110 supports more than one access control model, including a “Unix Perms” access control model (herein “Unix security style”) and an “NT ACL” access control model (herein “NT security style”).

Unix security style uses user IDs (UIDs) to identify users, and group IDs (GIDs) to identify groups to which those users belong. Unix security style associates the following access control limits with each file:

- o a UID for the owner;
- o a GID for the owner;
- o a set of “user” permissions—allowing permission to read, write, or execute the file by the owning user;
- o a set of “group” permissions—allowing permission to read, write, or execute the file by the owning user’s group; and

1
2 o a set of “other” permissions—allowing permission to read, write, or execute the
3 file by all other users.
4

5 Unix security style is supported by the NFS (“Network File System”) file
6 server protocol, possibly augmented with the NLM (“Network Lock Manager”) adjunct
7 file-locking protocol.
8

9 NFS is a stateless protocol, so each NFS file server request 121 includes the
10 UIDs and GIDs of the user making the request. The Unix client device 120 determines
11 the UIDs and GIDs for the user at a login time for the user, by reference to the system
12 password file (/etc/passwd) and group file (/etc/groups).
13

14 To enforce file access control using Unix security style, the file server 110
15 determines if the request 121 is from the owning user, from a user in the owning user’s
16 group, or from some other user. Responsive to this determination, the file server 110
17 uses one of the user permissions, the group permissions, or the other permissions, to de-
18 termine whether to allow the request 121.
19

20 NT security style uses security IDs (SIDs) to identify both users and
21 groups. NT security style associates the following access control limits with each file:
22

- 1 o an SID for the owner;
- 2
- 3 o an SID for the owner's group;
- 4
- 5 o an ACL (access control list).
- 6

7 The NT ACL includes one or more ACEs (access control entries), each of
8 which includes an SID indicating the user or group to which it applies, and a set of per-
9 missions. NT security style provides for the three Unix permissions (read, write, or exe-
10 cute), as well as "CHANGE PERMISSIONS" permission, "TAKE OWNERSHIP" per-
11 mission, "DELETE" permission, "DELETE CHILD" permission, and other permissions.

12
13 NT security style is supported by the CIFS ("Common Internet File Sys-
14 tem") protocol. NT security style is further described in the following articles: R. Rei-
15 chel, "Inside Windows NT Security", Windows/DOS Developers' Journal (April & May
16 1993), and in Stephen Sutton, "Windows NT Security Guide" (ISBN 0201419696).

17
18 CIFS is a session-based protocol, so the NT client device 120 transmits the
19 NT user name and password to the file server 110 at a session connect time, from which
20 the SIDs for the user and the user's groups are determined. The file server 110 can at-
21 tempt to authenticate the user itself, or (preferably) forward the NT user name and pass-
22 word to an NT primary domain controller.

1
2 To enforce file access control using NT security style, the file server 110
3 determines the SID for the user and the user's group, for the user making the request 121.
4 The file server 110 accumulates permissions granted to that user from ACEs that apply,
5 then subtracts permissions specifically denied to that user. Responsive to this accumula-
6 tion and subtraction, the file server 110 determines whether to allow the request 121.
7

8 Although a preferred embodiment of the invention is described with regard
9 to Unix security style and NT security style, the invention can readily be used with other
10 access control models, such as the "POSIX ACL" access control model supported by
11 some Unix devices, and by some other operating systems. The concepts and features of
12 the invention described herein can readily be used in a file server 110 supporting the
13 "POSIX ACL" access control model in addition to or instead of the access control models
14 described in detail herein, without further invention or undue experimentation. Accord-
15 ingly, the scope and spirit of the invention includes such file servers and methods for
16 their use.
17

18 The file server 110 designates each file 112 maintained in its file system
19 111 as having one specific access control model out of the plurality of access control
20 models it supports. In a preferred embodiment, each file 112 is designated to use either
21 Unix security style or NT security style. The file server 110 enforces the designated se-
22 curity style for each file 112 for all attempts to access that file 112. Thus, the file server

1 110 enforces the designated security style for all requests 121 made for that target file
2 112, whether those requests 121 come from Unix devices or NT devices, and whether
3 those requests 121 use the NFS file server protocol or the CIFS file server protocol.

5 *Access Control Enforcement*

7 If the file server 110 receives a request 121 for a target file 112, and the re-
8 quest 121 matches the security style target file 112, the file server 110 validates the re-
9 quest 121 against access control limits for that file 112 imposed by that security style.

11 The file server 110 thus recognizes and enforces at least the following cir-
12 cumstances:

- 14 o NT security style. The file 112 has NT security style and has a corresponding set
15 of access control limits (an NT ACL), which has been set by a client device 120
16 using the CIFS file server protocol.

18 If a client device 120 makes a request 121 to access the file 112 using the
19 CIFS file server protocol, the file server 110, if it is able to, enforces the NT ACL using
20 NT security style.

1 If the file server 110 is able to determine the NT user, either by communi-
 2 cation with an NT domain controller, or by reference to an NT user SID (security ID)
 3 database, the file server 110 is able to enforce the NT ACL using NT security style.

4
 5 If the file server 110 is not able to determine the NT user, it determines the
 6 equivalent Unix user, using a UID for a Unix user recorded for the CIFS file server pro-
 7 tocol session, and enforces the NT ACL as if the request 121 came from that Unix user.

8
 9 o Unix security style. The file 112 has Unix security style and has a corresponding
 10 set of access control limits (Unix Perms), which have been set by a client device
 11 120 using the NFS file server protocol

12
 13 If a client device 120 makes a request 121 to access the file 112 using the
 14 NFS file server protocol, the file server 110 enforces the Unix Perms using Unix security
 15 style.

16
 17 However, the file server 110 can also receive a request 121 that does not
 18 match the security style for the target file 112. The file server 110 can enforce the secu-
 19 rity style for the target file 112 against a non-matching client device 120 by validating
 20 either (1) a translated user ID for the client device 120 or (2) a translated set of access
 21 control limits for that file 112. As described herein, the file server 110 validates trans-
 22 lated user IDs for all Unix security style files 112, and preferably validates translated user

1 IDs for NT security style files 112 (when possible). Moreover, when the file server 110
2 validates translated access control limits for the file 112, the translated access control
3 limits are not recomputed for each request 121, but are cached with the file 112 for reuse.

4
5 The file server 110 thus also recognizes and enforces at least the following
6 circumstances:

- 7
8 o NT security style. The file 112 has NT security style and has a corresponding set
9 of access control limits (an NT ACL), which has been set by a client device 120
10 using the CIFS file server protocol.

11
12 If a client device 120 makes a request 121 to access the file 112 using the
13 NFS file server protocol, the file server 110 determines the Unix user, associated with the
14 client device 120, that is making the request 121. The Unix user has a UID (user ID).

15
16 In a preferred embodiment, the file server 110 maps the Unix user into an
17 equivalent NT user. The file server 110 translates the UID into an SID (security ID) that
18 is an equivalent user to the Unix user. The file server 110 enforces the access control
19 limits (the NT ACL) for the equivalent NT user (the SID).

20
21 The file server 110 performs the following process to map the Unix user
22 into an equivalent NT user:

- 1
- 2 o The client device 120 contacts the file server 110 using the NFS file server proto-
- 3 col. The NFS file server protocol request 121 includes a UID for the Unix user as-
- 4 sociated with the client device 120.
- 5
- 6 o The file server 110 looks up the UID in the Unix password file (/etc/passwd), and
- 7 thus identifies the Unix user name for the Unix user. The Unix user name is an al-
- 8 phanumeric string identifying the Unix user.
- 9
- 10 o The file server 110 translates the Unix user name into an NT user name using a
- 11 selected mapping file. Similar to the Unix user name, the NT user name is an al-
- 12 phanumeric string identifying the NT user. If there is no such translation for the
- 13 Unix user name, the file server 110 uses the Unix user name, without translation,
- 14 as the NT user name.
- 15
- 16 In a preferred embodiment, the mapping file includes a set of records each identi-
- 17 fying an NT user by NT user name, and associating an equivalent Unix user name
- 18 with the NT user name.
- 19
- 20 o The file server 110 contacts an NT domain controller to determine an SID for the
- 21 NT user name. If there is no such NT user, the file server 110 uses a selected pa-

parameter for unmapped Unix users. In a preferred embodiment, this selected parameter is set to the NT user "guest" by default.

o The file server 110 contacts the NT domain controller to obtain the SIDs of all groups for which the NT user is a member.

The file server 110 caches UID-to-SID mappings for a period of time, preferably about a few hours.

In an alternative preferred embodiment, or if the file server 110 is unable to map Unix users to NT users (for example, if domain authentication has been turned off), the file server 110 maps the NT ACL into a no less restrictive set of Unix Perms. The file server 110 determines these Unix Perms in response to the NT ACL and in response to the Unix user. The file server 110 enforces the mapped access control limits (the Unix Perms) for the actual Unix user (the UID).

The file server 110 might perform dynamic permission mapping, in which the file server 110 maps the NT ACL into a set of Unix Perms at the time the mapping is required. In a preferred embodiment, the file server 110 caches the translated Unix Perms with the file 112. Accordingly, for validating access control limits, the file server 110 performs static permission mapping, in which the file server 110 maps the NT ACL into a set of Unix Perms at the time the NT ACL is set.

1
2 The file server 110 performs the following process to achieve static permis-
3 sion mapping:

4
5 o The file server 110 determines the NT user that is the owner of the file 112, and
6 maps the NT user into an equivalent Unix user (the file server 110 maps the SID
7 for the NT user into a UID for a Unix user).

8
9 o The file server 110 examines the NT ACL for the file 112 and determines whether
10 there are any “deny access” provisions.

11
12 o If the NT ACL for the file 112 has no “deny access” provisions, the file server 110
13 generates a set of Unix Perms having an entry for “user permissions,” consistent
14 with the file access control limits provided by the NT ACL. The file server 110
15 sets the Unix Perms for “group permissions” equal to the Unix Perms for “other
16 permissions.” The file server 110 sets the Unix Perms for “other permissions”
17 equal to the NT ACL entry for “everyone,” if one exists.

18
19 o If the NT ACL for the file 112 does have “deny access” provisions, the file server
20 110 rejects the request 121.

1 Because static permission mapping is not responsive to the particular user
2 making the request 121, the file server 110 does not attempt to determine what the provi-
3 sions of the NT ACL are for that particular user.

4
5 o Unix security style. The file 112 has Unix security style and has a corresponding
6 set of access control limits (Unix Perms), which have been set by a client device
7 120 using the NFS file server protocol

8
9 If a client device 120 makes a request 121 to access the file 112 using the
10 CIFS file server protocol, the file server 110 determines the NT user, associated with the
11 client device 120, that is making the request 121. The NT user has an SID (session ID).

12
13 The file server 110 maps the NT user into an equivalent Unix user. The file
14 server 110 translates the SID into a UID that is an equivalent user to the NT user. The
15 file server 110 enforces the access control limits (the Unix Perms) for the equivalent Unix
16 user (the UID).

17
18 The file server 110 performs the following process to map the NT user into
19 an equivalent Unix user:

- 1 o The client device 120 starts a CIFS session (the client device 120 first contacts the
2 file server 110 using the CIFS file server protocol). The client device 120 trans-
3 mits its NT user name to the file server 110.
4
- 5 o The file server 110 translates the NT user name into a Unix user name using a
6 mapping file. If there is no such translation for the NT user name, the file server
7 110 uses the NT user name, without translation, as the Unix user name.
8
- 9 o The file server 110 looks up the Unix user name in the Unix password file
10 (/etc/passwd), and thus identifies the Unix user, the UID for the Unix user, the
11 Unix user's primary group, and the primary GID (group ID) for the Unix user. If
12 there is no such Unix user name in the Unix password file, the file server 110 uses
13 a selected parameter for unmapped NT users. In a preferred embodiment, this se-
14 lected parameter is set to the Unix user "nobody" by default.
15
- 16 o The file server 110 looks up the Unix user name in the Unix group file
17 (/etc/groups) to determine any other groups and any other GIDs associated with
18 the Unix user.
19

20 *Reading and Modifying Access Control Limits*

21

1 Each file 112 has its security style set by the file server 110 so that either
2 (a) a request 121 to perform an operation on the file 112, or (b) a request 121 to perform
3 an operation that sets the access control limits for the file 112, produce expected results.
4

5 When the file 112 is first created, the file server 110 sets the security style
6 for the file 112 equal to a security style associated with the file server protocol used to
7 create it. (This is limited by restrictions imposed by access control trees, described
8 herein.) Thus, if the file 112 is created using the NFS file server protocol, the security
9 style for the file 112 is set to Unix security style. Similarly, if the file 112 is created us-
10 ing the CIFS file server protocol, the security style for the file 112 is set to NT security
11 style.
12

13 When the file 112 has its access control limits modified, the file server 110
14 sets the security style for the file 112 equal to a security style associated with the new ac-
15 cess control limits. (This is limited by restrictions imposed by access control trees, de-
16 scribed herein.) Thus, if a client device 120 sets a set of Unix Perms for the file 112, the
17 security style for the file 112 is set to Unix security style. Similarly, if a client device 120
18 sets an NT ACL for the file 112, the security style for the file 112 is set to NT security
19 style.
20

21 The file server 110 can receive a request 121 to read or view the access
22 control limits for a file 112. Also, when the file server 110 receives a request 121 to

1 make an incremental change to the access control limits for a file 112, it determines the
2 current access control limits for the file 112 before making the incremental change.

3
4 The file server 110 thus recognizes and enforces at least the following cir-
5 cumstances:

- 6
7 o NT security style. The file 112 has NT security style and has a corresponding set
8 of access control limits (an NT ACL), which has been set by a client device 120
9 using the CIFS file server protocol.

10
11 If a client device 120 makes a request 121 to read or modify the access
12 control limits for the file 112 using the NFS file server protocol, the file server 110 de-
13 termines the Unix user, associated with the client device 120, that is making the request
14 121.

15
16 The file server 110 performs the same process to map an NT ACL into a set
17 of Unix Perms as described above for validation of file access control, with the following
18 exceptions:

19
20 Unlike validation of access control limits, the file server 110 treats transla-
21 tion of access control limits differently for requests 121 to read or modify the access
22 control limits for the file 112.

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Preferably, the file server 110 performs dynamic permission mapping, in which the file server 110 maps the NT ACL into a set of Unix Perms at the time the mapping is required. NT security style is richer than Unix security style—for example, Unix security style has no “deny access” provisions. Thus, it is possible for the file server 110 to map the NT ACL into a set of Unix Perms that appears different for different Unix users. For example, if the NT ACL, for a file 112 whose owner is Charles, specifically provides read access to Allen but specifically denies read access to Beth, the file server 110 will provide different Unix perms to each of the users Allen and Beth. One set will allow read access by Allen’s group and one set will disallow read access by Beth’s group, in harmony with the access provided by the actual NT ACL.

The file server 110 performs the following process to achieve dynamic permission mapping:

- o The file server 110 determines the NT user that is the owner of the file 112, and maps the NT user into an equivalent Unix user (the file server 110 maps the SID for the NT user into a UID for a Unix user).
- o The file server 110 examines the NT ACL for the file 112 and determines whether there are any “deny access” provisions.

1 o If the NT ACL for the file 112 has no "deny access" provisions, the file server 110
2 generates a set of Unix Perms having entries for "user permissions" and "other
3 permissions," consistent with the file access control limits provided by the NT
4 ACL. The file server 110 sets the Unix Perms for "group permissions" equal to
5 the Unix Perms for "other permissions."

6
7 o If the NT ACL for the file 112 does have "deny access" provisions, the file server
8 110 attempts to determine if any apply to the particular Unix user. If the file
9 server can tell, it generates a set of Unix Perms that reflect the access control lim-
10 its currently available for this particular file 112 and this particular Unix user. If
11 the file server 110 cannot tell, it rejects the request 121. (Alternatively, the file
12 server 110 could reject the request 121 if there are any "deny access" provisions in
13 the NT ACL.)

14
15 In alternative embodiments, the file server 110 may perform static permis-
16 sion checking, similar to validation of access control limits, for requests 121 to read or
17 modify the access control limits for the file 112.

18
19 If the request 121 attempts to modify attributes of the file 112 that have no
20 effect on access control limits for the file 112 (such as access time or modify time), the
21 file server 110 makes those modifications without change to the access control limits for
22 the file 112.

1
2 If the request 121 attempts to modify some but not all access control limits
3 for the file 112, the file server 110 generates a set of Unix Perms in response to the NT
4 ACL for the file 112, as described above. The file server 110 modifies the generated
5 Unix Perms as specified by the request 121. If the file server 110 cannot generate a set of
6 Unix Perms in response to the NT ACL for the file 112, the file server 110 rejects the re-
7 quest 121.

8
9 One difference in setting access control limits is that, according to NT secu-
10 rity style, files 112 can be specifically set to be "READ-ONLY." According to Unix se-
11 curity style, files are set to be read only by clearing the WRITE permission for the owner
12 of the file 112. When a client device 120 using the CIFS file server protocol attempts to
13 set the READ-ONLY attribute of a file 112 with Unix security style, the file server 110
14 clears the WRITE permission for the owner of the file 112 in the Unix Perms for that file
15 112.

16
17 o Unix security style. The file 112 has Unix security style and has a corresponding
18 set of access control limits (Unix Perms), which have been set by a client device
19 120 using the NFS file server protocol

20

1 The file server 110 performs the following process to map a set of Unix
2 Perms into an NT ACL for display or modification of those Unix Perms by a CIFS client
3 device 120:

- 4
- 5 o The file server 110 generates an NT ACL entry for "owner," providing the same
6 access control limits as the Unix Perms entry for "user permissions."
- 7
- 8 o The file server 110 generates an NT ACL entry for "everyone," providing the
9 same access control limits as the Unix Perms entry for "other permissions."
- 10
- 11 o If possible, the file server 110 generates an NT ACL entry for the actual request-
12 ing user, providing the same access control limits as the Unix Perms entry for that
13 user. This step could require mapping the Unix user into an equivalent NT user
14 using the UID-to-SID cache.
- 15

16 Similar to modification of an NT ACL entry by a Unix user, if the request
17 121 (for modification of Unix Perms by an NT user) attempts to modify attributes of the
18 file 112 that have no effect on access control limits for the file 112, the file server 110
19 makes those modifications without change to the access control limits for the file 112.

20

21 If the request 121 attempts to modify some but not all access control limits
22 for the file 112, the file server 110 generates an NT ACL in response to the set of Unix

1 Perms for the file 112, as described above. The file server 110 modifies the generated
2 NT ACL as specified by the request 121.

4 *Access Control Subtrees*

6 In a preferred embodiment, the files 112 in the file system 111 are organ-
7 ized into a tree, having a set of branch nodes and a set of leaf nodes. One branch node of
8 the tree is a root node, and each branch node of the tree is a root node for a subtree of the
9 tree. In the file system 111, each branch node is a directory, and each leaf node is a file
10 112. A directory is a type of file 112 that includes information about those branch nodes
11 and leaf nodes in a subtree for which it is the root node.

13 The file server 110 associates a limited set of access control models with
14 each subtree. In a preferred embodiment in which the file server 110 supports Unix secu-
15 rity style and NT security style, the file server 110 designates each subtree as being NT-
16 only format, Unix-only format, or mixed format.

18 NT-only Format

20 When the file server 110 designates a subtree as being NT-only format, it
21 restricts creation of files 112 within that subtree to files 112 having NT security style.

1 The file server 110 also prohibits changing the access control model of files 112 within
2 that subtree to other than NT security style.

3
4 According to NT security style, new files 112 inherit NT ACL settings
5 from their parent nodes. If a client device 120 using the NFS file server protocol at-
6 tempts to create a file 112 in a subtree having NT-only format, that file 112 can only be
7 created by the Unix user equivalent to the NT user who is the NT-owner of the root node
8 of the subtree. The file server 110 determines if the Unix user making the request 121 is
9 the equivalent by (a) mapping the SID for the NT user who is the owner into an equiva-
10 lent UID; (b) storing that UID in its record for the file 112; and (c) comparing that UID
11 with the UID in the request 121.

12
13 According to NT security style, there is a particular "DELETE" permission
14 and a particular "DELETE-CHILD" permission. If the file server 110 is unable to deter-
15 mine if a Unix user has these permissions, it rejects requests 121 to delete files 112 in
16 NT-only format subtrees, unless the request 121 is from the owner of the file 112 (the
17 equivalent Unix user of the NT user who is the owner) or the Unix user "root".

18
19 According to NT security style, there is a particular "CHANGE-
20 PERMISSION" permission and a particular "TAKE-OWNER" permission. If the file
21 server 110 is unable to determine if a Unix user has these permissions, it denies requests
22 121 to set any permissions for files 112 in a NT-only format subtree, unless the request

1 121 is from the owner of the file 112 (the equivalent Unix user of the NT user who is the
2 owner) or the Unix user "root".

3 4 Unix-only Format

5
6 Similarly, when the file server 110 designates a subtree as being Unix-only
7 format, it restricts creation of files 111 within that subtree to files 111 having Unix secu-
8 rity style. The file server 110 also prohibits changing the access control model of files
9 111 within that subtree to other than Unix security style. Attempts to set an NT ACL
10 would change the access control model for that file 112 to NT security style, and so are
11 rejected in a Unix-only format subtree.

12
13 When a client device 120 using the CIFS file server protocol creates a file
14 112 in a Unix-only format subtree, the file server 110 sets the owner of the file 112 to the
15 Unix user equivalent to the NT user making the request 121. The file server 110 maps
16 the SID for the NT user to a UID for an equivalent Unix user, and uses that UID to set the
17 owner of the file 112.

18
19 According to Unix security style, there is no "CHANGE-PERMISSION"
20 permission or "TAKE-OWNER" permission. The file server 110 always denies requests
21 121 to set these permissions for files 112 in a Unix-only format subtree.

1 Mixed Format

2

3 When the file server 110 designates a subtree as being mixed format, it al-
4 lows creation of files 111 with either Unix security style or NT security style. The file
5 server 110 does not prohibit changing the access control model of files 111 within that
6 subtree to either Unix security style or NT security style.

7

8 An administrator for the file server 110 can change the designation of a
9 subtree from a first format to a second format (for example, from mixed format to either
10 NT-only format or Unix-only format). When the second format is possibly incompatible
11 with the first format (for example, a subtree changed to NT-only format includes nodes
12 that are Unix security style), the file server 110 converts those files 112 with incompati-
13 ble access control models as it sets permissions for those files 112. Requests 121 for a
14 file 112 which only check permissions are still validated using the access control model
15 in place for the file 112.

16

17 Although the invention is described herein with regard to only two access
18 control models, the invention can readily be used with three or more access control mod-
19 els. In such alternative embodiments, there are a larger number of possible subtree for-
20 mats, including subtree formats that restrict the files 112 within that subtree to one of a
21 set of multiple access control models, less than the set of all access control models recog-
22 nized by the file server 110.

1
2 In a preferred embodiment, the root node of the file system 111 is desig-
3 nated as having mixed format. Client devices 120 that are owners of a subtree can mod-
4 ify the format of a subtree by request 121 to the file server 110; thus, client devices 120
5 can modify subtrees to have NT-only format, Unix-only format, mixed format. When a
6 new subtree is created, the file server 110 designates the new subtree as having the same
7 format as its parent; thus, mixed format if it is created within a subtree that is already
8 mixed format (the default), NT-only format if it is created within a subtree that is already
9 NT-only format, and Unix-only format if it is created within a subtree that is already
10 Unix-only format.

11 12 *Alternative Embodiments*

13
14 Although preferred embodiments are disclosed herein, many variations are
15 possible which remain within the concept, scope, and spirit of the invention, and these
16 variations would become clear to those skilled in the art after perusal of this application.